SECURE VAULT SCHEME IN THE CLOUD OPERATING MODEL

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- Data Privacy Vaults
 - What do they do?
 - Who's making them? Who's using them?
- Our framework
- Building Block: Tokenization Scheme
 - Definition
 - Security Notions
- Complete Vault Scheme Construction
 - Construction

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DATA-PRIVACY VAULTS

- → What do they do?
- → Who is using them?



Data-Privacy Vault

A literal vault (safe) of information -



- Users can store their private information
- Generate access passes
 - give access to selected parties
- Data hidden from others



Data-Privacy Vault

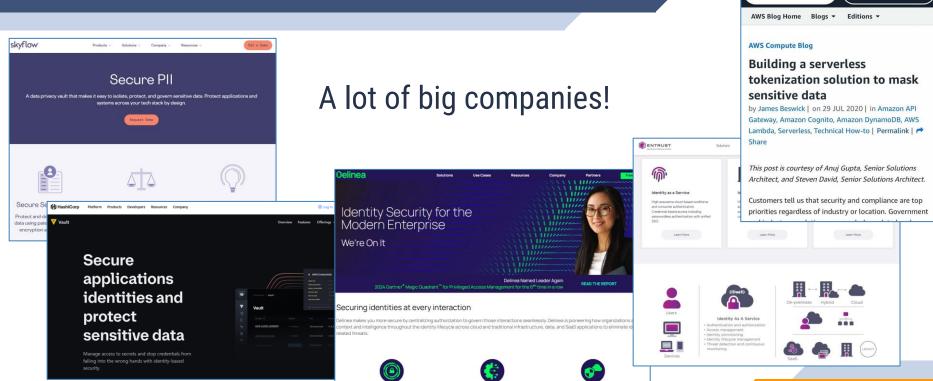
Access Levels

- Different levels of access passes
 - provides access to different data
- Eg. Level 1: Non-sensitive data
 Level 2: Both non-sensitive and
 sensitive data





> Who's making & using them?



We make you

more productive

We future proof

identity security

We make you

more secure

aws

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Who's making & using them?

A lot of big companies!

increasingly being used for storing information used for training LLM models

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OUR FRAMEWORK



Our Framework

3 parties:



TRUSTED SERVER



USER





ANALYST

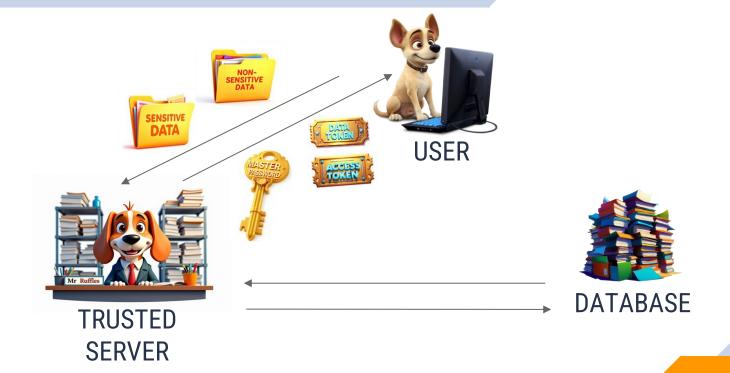


Our Framework - Functionality

- STORE -
 - User stores the information and gets back access passes
- ACCESS -
 - Level 1 access to non-sensitive information
- **RETRIEVE** -
 - Level 2 access to all the information

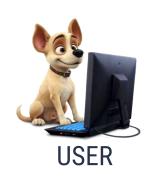


STORE





RETRIEVE



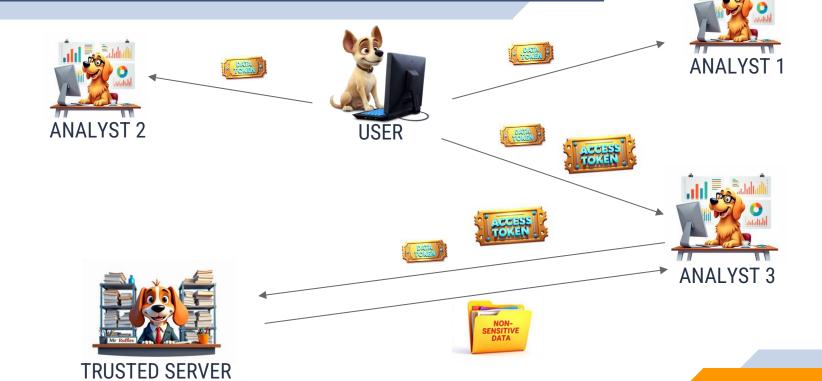




TRUSTED SERVER



ACCESS





Non-Sensitive Information:

No access Token -> No non-sensitive information

Sensitive Information:

No master password -> No sensitive information

(even with access token)



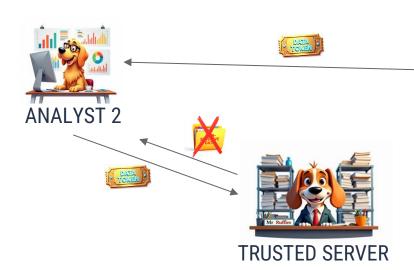
IND-CATA





Non-Sensitive Information:

No access Token -> No non-sensitive information





USER



Non-Sensitive Information:

No access Token -> No non-sensitive information

```
Game IND-CATA
 1: b \leftarrow \$ \{0,1\}
 2: PP \leftarrow s \operatorname{Param}(1^{\lambda})
       (state, m_0 = (m_{s,0}, m_{ns,0}), m_1 = (m_{s,1}, m_{ns,1}))
                                     \leftarrow \mathcal{A}^{\mathsf{Store}(PP,.),\mathsf{Access}(PP,.)}(PP,1^{\lambda})
         if m_{ns,0} = m_{ns,1}:
            return 0
 5:
         (m\text{-}tok^*, data\text{-}tok^*) = \mathsf{Store}(\bot, PP, m_b)
         access-tok^* = \mathcal{H}(10, m-tok^*)
        b' \leftarrow \mathcal{A}^{\mathsf{Store}(PP,.),\mathsf{Access}(PP.)}(PP,access\text{-}tok^*,state)
 9: return (b=b')
```







Sensitive Information:

No master password -> No sensitive information (even with access token)





Game IND-SIA

```
1: b \leftarrow s \{0,1\}
2: PP \leftarrow s \operatorname{Param}(1^{\lambda})
       (state, m_0 = (m_{s,0}, m_{ns,0}), m_1 = (m_{s,1}, m_{ns,1}))
                                  \leftarrow \mathcal{A}^{\mathsf{Store}(PP,.),\mathsf{Access}(PP,.)}(PP,1^{\lambda},DB)
       ((m-tok^*, data-tok^*), DB') = Store(\bot, PP, m_{ns,b}, DB)
5: b' \leftarrow A^{\text{Store}(PP,.),\text{Access}(PP,.)}(PP, state, DB')
6: return (b = b')
```

Sensitive Information:

No master password -> No sensitive information (even with access token)



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BUILDING BLOCK: TOKENIZATION SCHEME



Tokenization Scheme - Model

- Cloud Operating Model
 - Database may get compromised
- LLM-training friendly
 - Must be deterministic (same token for same data)
- The user should not have to store a secret key



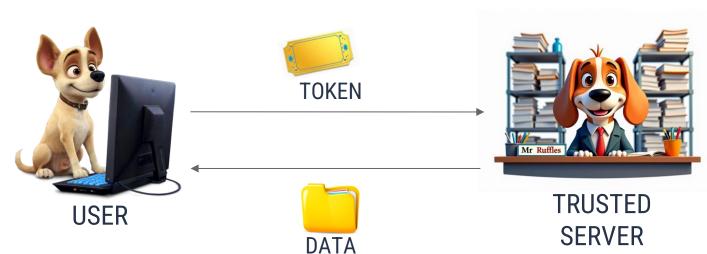
Tokenization Scheme







Tokenization Scheme







Tokenization Scheme

- **Tokenization Procedure:**
 - ► Takes message m as input
 - Generates a token tok and adds an entry to the database
 - ► Returns tok
- **DeTokenization Procedure:**
 - Takes token tok as input
 - If it exists in database, returns corresponding message m, else returns \perp .





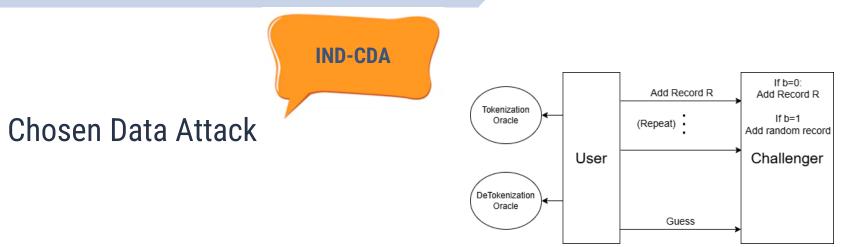
Chosen Data Attack



Chosen Distribution Database Attack







Chosen Distribution Database Attack





Our Construction

```
\mathsf{Tok}(m,\overline{m},k_1) // for basic construction k_1=\perp, m=\overline{m}
                                                                                    DeTok(tok)
 1: if k_1 = \perp
                                                                                      1: if \exists (c_1||t_1||c_2||t_2) \in DB \ s.t. \ t_2 = \mathcal{H}_4(tok)
           k_1 \leftarrow \mathcal{H}_1(m)
                                                                                                 m \leftarrow \mathsf{Dec}\left(\mathcal{H}_3(tok), c_2\right)
 3: if \exists (c_1||t_1||c_2||t_2) \in DB \ s.t. \ t_1 = \mathcal{H}_2(m)
                                                                                                 return m
           tok \leftarrow Dec(k_1, c_1)
                                                                                             else return
           return tok
 6: tok \leftarrow \$ \{0,1\}^{\tau}
                                                    token generated randomly,
 7: c_1 \leftarrow \mathsf{Enc}(k_1, tok)
 8: t_1 \leftarrow \mathcal{H}_2(m)
        k_2 \leftarrow \mathcal{H}_3(tok)
                                                   c_1 = \operatorname{Enc}(H(m), tok), t_1 = H'(m)

c_2 = \operatorname{Enc}(H(tok), m), t_2 = H'(tok)
        c_2 \leftarrow \mathsf{Enc}(k_2, \overline{m})
      t_2 \leftarrow \mathcal{H}_4(tok)
12: d = c_1||t_1||c_2||t_2|
13: \quad DB = DB \cup \{d\}
                                                    store {c1, t1, c2, t2} as one record
14: return tok
```



Our Construction

```
\mathsf{Tok}(m,\overline{m},k_1) // for basic construction k_1 = \perp, m = \overline{m}
                                                                                               DeTok(tok)
 1: if k_1 = \perp
                                                                                                 1: if \exists (c_1||t_1||c_2||t_2) \in DB \ s.t. \ t_2 = \mathcal{H}_4(tok)
             k_1 \leftarrow \mathcal{H}_1(m)
                                                                                                             m \leftarrow \mathsf{Dec}\left(\mathcal{H}_3(tok), c_2\right)
         if \exists (c_1||t_1||c_2||t_2) \in DB \ s.t. \ t_1 = \mathcal{H}_2(m)
                                                                                                             return m
             tok \leftarrow Dec(k_1, c_1)
                                                                                                         else return
             return tok
         tok \leftarrow \$ \{0,1\}^{\tau}
         c_1 \leftarrow \mathsf{Enc}(k_1, tok)
 8: t_1 \leftarrow \mathcal{H}_2(m)
         k_2 \leftarrow \mathcal{H}_3(tok)
         c_2 \leftarrow \mathsf{Enc}(k_2, \overline{m})
        t_2 \leftarrow \mathcal{H}_4(tok)
12: d = c_1||t_1||c_2||t_2
         DB = DB \cup \{d\}
14: return tok
```

```
if a record exists with
m = Dec(H(tok), c_2)
```

Our construction is both IND-CDA and IND-CDDA secure

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FULL VAULT SCHEME CONSTRUCTION



> Vault Scheme Construction

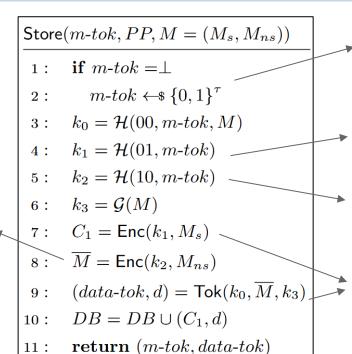
We use the Tokenization Scheme as a building block

For simplicity of presentation, assume that the message space is same as the token space.



Vault Scheme Construction - STORE

n-s data is encrypted with the access token as key



if token does not exist beforehand, it is generated.

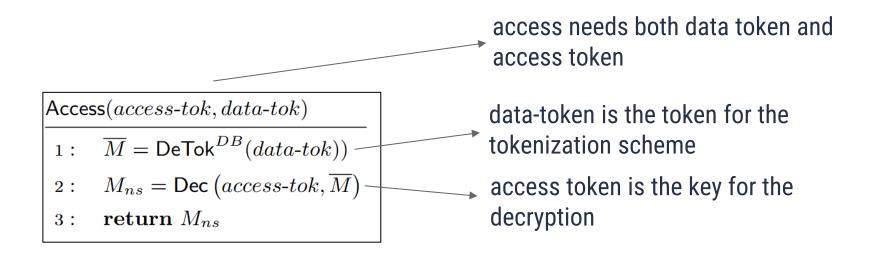
two hashes of the master token, one for sensitive, one for non-sensitive data

the one for the non-sensitive data is used as the access token

the non-sensitive data ciphertext is tokenized and that token is used as the data-token



Vault Scheme Construction - ACCESS





Vault Scheme Construction - RETRIEVE

Retrieve(m-tok, data-tok)

1:
$$((C_1,d),\overline{M}) = \mathsf{DeTok}^{DB}(data-tok))$$

$$2: \quad k_1 = \mathcal{H}(0, m\text{-}tok)$$

$$3: k_2 = \mathcal{H}(1, m\text{-}tok)$$

$$4: \quad M_{ns} = \mathsf{Dec}\left(k_2, \overline{M}\right)$$

$$5: \quad M_s = \mathsf{Dec}\left(k_1, C_1\right)$$

6: **return**
$$M = (M_s, M_{ns})$$

else return

retrieve needs both master password and access token

data-token is the token for the tokenization scheme

derive the access token and use it to access the non-sensitive data

use master password to decrypt sensitive data

Security

- The IND-CATA of the vault construction follows from -
 - Randomness of the hash function
 - IND-CDA and IND-CDDA security of the underlying Tokenization scheme.
- The IND-SIA follows from
 - randomness of the hash functions
 - IND-CPA security of encryption scheme
 - IND-CDDA of tokenization scheme



Thank You!